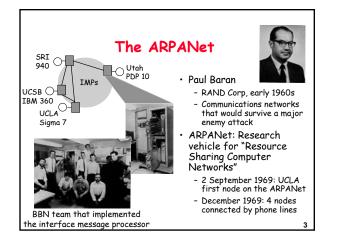
CENG 577 Advanced Services in Communications

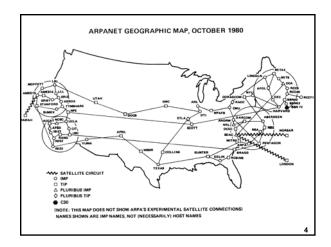
Internet Architecture and Advanced Services in Converged Networks

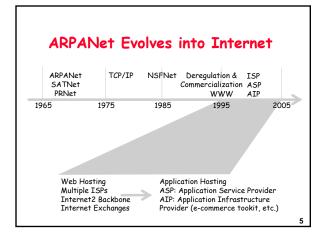
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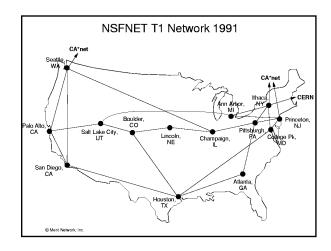
Outline

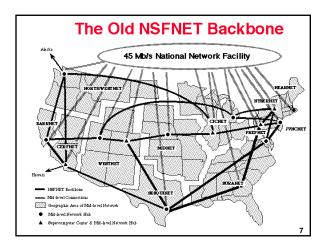
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- Principles of Data Communications
- Higher Internet View and New Trends
- Business Trends
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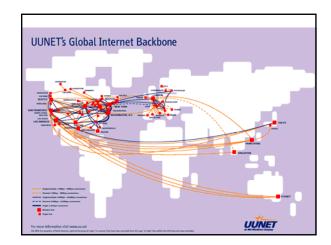


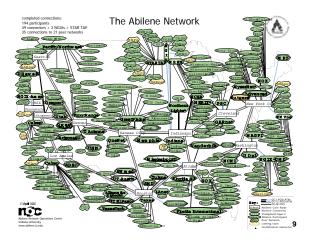


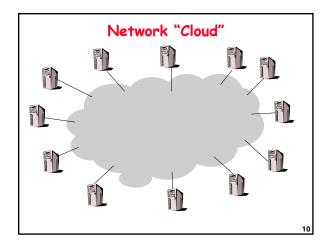


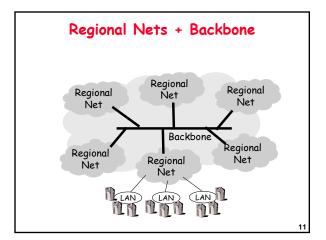


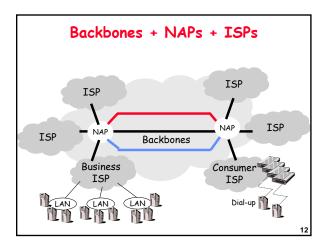


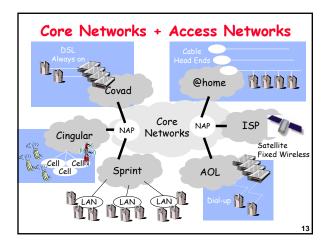


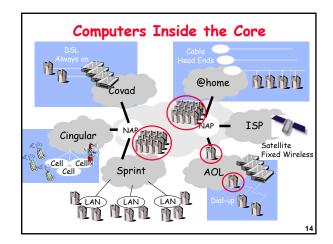


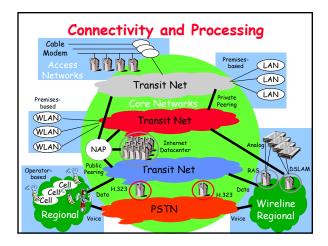


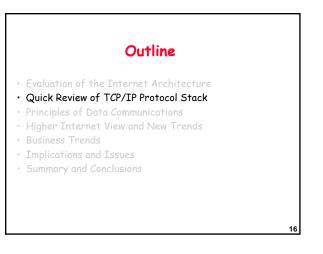


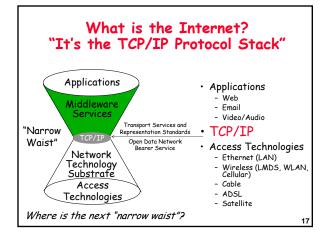


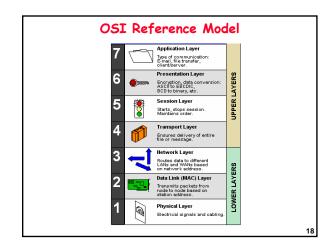


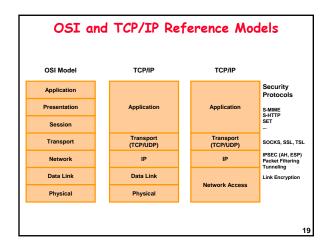


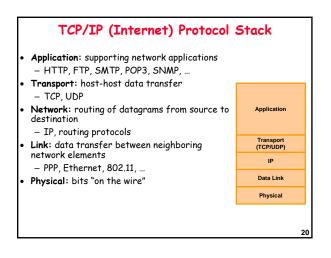


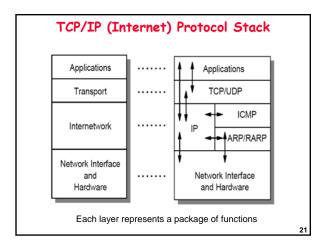


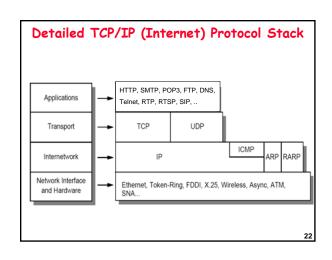


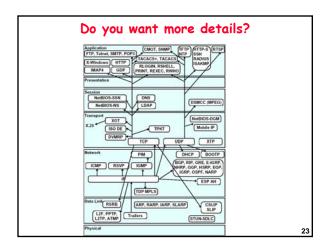


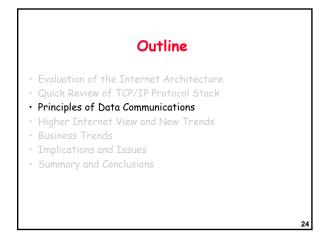


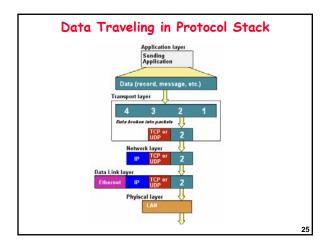


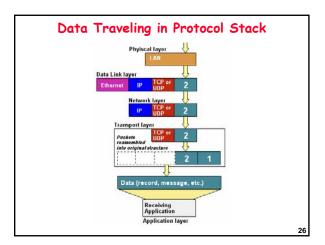










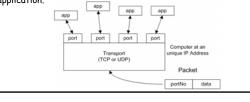


Connection in TCP/IP

- A connection between two machines in TCP/IP is defined by:
- Transport layer protocol (TCP or UDP)
- IP address of local machine
- Port number used on the local machine
- IP address of remote machine
- Port number used on the remote machine

Transport-Level Protocols: Ports

- TCP and UDP protocols use **ports** to map incoming data to a particular process running on a computer.
- IP Datagram identifies the **host** and the **port** that it's destined for.
- The computer is identified by its 32-bit IP address, which IP uses to deliver data to the right computer on the network.
- Ports are identified by a 16-bit integer number, ranging from 0 to 65535, which TCP/UDP use to deliver the data to the right application.



Well-known Ports Port numbers between 0 and 1023 are restricted (well-known ports) -- they are reserved for use by well-known services such as HTTP and FTP and other system services. Your applications should not attempt to bind to these ports.

Protocol	Port	Encoding	Purpose	
Echo	7	TCP/UDP	Test protocol used to verify that two machines are able to connect by having one echo back the other's input.	
Discard	9	TCP/UDP	Less useful test protocol that ignores all data received by the server.	
Daytime	13	TCP/UDP	Provides an ASCII representation of the current time on the server.	
ftp-data	20	TCP	FTP uses two well-known ports. This port is used to transfer files.	
FTP	21	TCP	This port is used to sent ftp commands like, "put" and "get".	
TELNET	23	TCP	A protocol used for interactive, remote command-line sessions.	
SMTP	25	TCP	"Simple Mail Transfer Protocol" is used to send email between machines.	
Time	37	TCP/UDP	A time server returns the number of seconds that have elapsed on the host machine since midnight, January 1, 1900, as a four-byte, signed, big-endian int	
Whois	43	TCP	Simple directory service for Internet network administrators,	
Finger	79	TCP	It gets information about a user or users.	
нттр	80	TCP	Hyper Text Protocol is the underlying protocol of the World Wide Web.	
POP3	110	TCP	Post Office Protocol version 3 is a protocol for the transfer of accumulated form the host to sporadically connected clients.	
NNTP	119	TCP	Usenet news transfer. More formally known as the "Network News Transfer Protocol".	
SNMP	161/16 2	UDP	Simple Network Managment Protocol is used in management of TCP/IP.	
RMI Reg.	1099	TCP	The RMI Registry is a registry service for Java remote objects.	
Servlets	8080	TCP	Java Server API and Servlets is a web server from Sun that runs on port 8080 by default, not port 80.	

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Transport-Level Protocols: TCP

 TCP is a reliable and connection-oriented communication protocol on top of the unreliable, unsequenced functionality of IP.

Reliable:

- TCP provides extensive error-checking capabilities.
- TCP provides ${\bf reliable\ stream\ delivery}.$ This reliable stream delivery ensures that the data arrives in the same sequence and state in which it was sent.

Connection-oriented:

- The TCP system relies on a **virtual circuit** that is established between the requesting machine and its target.
- This circuit is opened via a 3-part process, often referred to as the 3-part handshake.

Transport-Level Protocols: TCP

- Because of the reliable and sequenced nature of TCP sockets, they often are called stream sockets; you can read and write data in continuous streams of bytes without worrying about packets, headers, and so on.
- TCP is the chief protocol employed on the Internet.
- It facilitates such mission-critical tasks as file transfers and remote sessions.
- Stream socket functionality in Java is provided by the classes java.net.ServerSocket and java.net.Socket.

Transport-Level Protocols: UDP

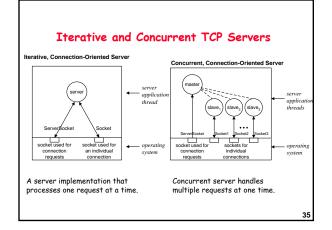
UDP is an unreliable and connectionless communication protocol.

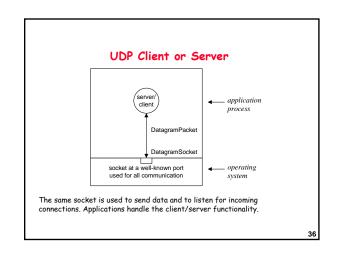
- Datagram-based communication.
- Datagram packets are prepared by the applications.
- IP Address + Port Number are put into datagram.
- UDP-based communication is like sending letters to a post office. Not reliable but fast compared to TCP.
- Datagram socket functionality in Java is provided by the classes java.net.DatagramSocket and java.net.DatagramPacket.

TCP-Based Communication Steps

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- A server application opens a socket to establish a connection with another application (client) by binding a socket to a port number. (registering the application with the system to receive all data destined for that port.)
- Server: TCP Socket = Port Number (Well-known)
- Client: TCP Socket = IP Address + Port Number (server's port)
- When a **client** makes a request from the server's port, input and/or output streams are created on the socket depending on the protocol used between the server and the client.
- No two applications can bind to the same port: Attempts to bind to a port that is already in use will fail.
- Stream based (like a phone call)
- Uses 3-way handshake, reliable but slow (compared to UDP)





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Internet Applications

- Variations on three themes
- distinguish protocol vs. application behavior
- Messaging
- datagram model \rightarrow no direct confirmation of final receipt
- email (optional confirmation now) and IM
 emphasis on interoperation (SMS, pagers, ...)
- delays measured in minutes

• Retrieval & query (request/response)

- "client-server"
- ftp, HTTP
- RPC (Sun RPC, DCE, DCOM, Corba, XML-RPC, SOAP)

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- emphasis on fast & reliable transmission
- delays measured in seconds

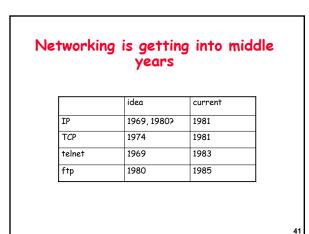
Internet applications, cont'd

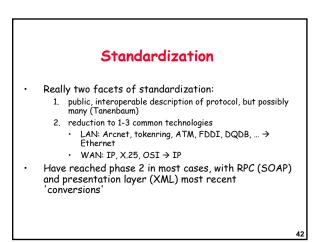
• Continuous Media

- generation rate ~ delivery rate ~ rendering rate
- audio, video, measurements, control
 - » Internet telephony
 - » Multimedia conferencing
- related: streaming media \rightarrow slightly longer timescales for rate matching
- » video-on-demand
- emphasis is on *timely* and low-loss delivery \rightarrow *real-time*
- delays measured in milliseconds
- focus of this course

Internet protocols

- Protocols support these applications:
 - data delivery
 » HTTP, ftp data part, SMTP, IMAP, POP, NFS, SMB, RTP
 identifier mapping (id → id, id → data)
 - » ARP, DNS, LDAP
 - configuration (= specialized version of identifier → data)
 » DHCP, ACAP, SLP, NETCONF, SNMP
 - control and setup
 - » RTSP, SIP, ftp control, RSVP, SNMP, BGP and routing protocols
- May be integrated into one protocol or general service function ("middleware"?)





Technologies at ~30 years

- Other technologies at similar maturity level:
 - air planes: 1903 1938 (Stratoliner)
 - cars: 1876 1908 (Model T)
 - analog telephones: 1876 1915 (transcontinental telephone)
 - railroad: 1800s -- ?

Observations on progress

- 1960s: military → professional → consumer
 now, often reversed
- Oscillate: convergence \rightarrow divergence
- continued convergence clearly at physical layer
- niches larger → support separate networks
- Communications technologies rarely disappear (as long as operational cost is low):
- exceptions
 - » telex, telegram, semaphores → fax, email
 - » X.25 + OSI, X.400 → IP, SMTP
 - analog cell phones

History of networking

- History of networking = non-network applications migrate
 - postal & intracompany mail, fax → email, IM
 - broadcast: TV, radio
 - interactive voice/video communication \rightarrow VoIP
 - information access \rightarrow web, P2P
 - disk access → iSCSI, Fiberchannel-over-IP

Network evolution

- Only three modes, now thoroughly explored:
 - packet/cell-based
 - message-based (application data units)
 - session-based (circuits)
- Replace specialized networks
 - left to do: embedded systems
 - » need cost(CPU + network) < \$10
 - » cars
 - » industrial (manufacturing) control
 - » commercial buildings (lighting, HVAC, security; now LONworks)

 - » remote controls, light switches
 - » keys replaced by biometrics

New applications

- · New bandwidth-intensive applications - Reality-based networking

 - (security) cameras
- Distributed games often require only low-bandwidth control information
- current game traffic ~ VoIP
- · Computation vs. storage vs. communications - communications cost has decreased less rapidly than storage costs

Security challenges

- · DOS, security attacks \rightarrow permissions-based communications - only allow modest rates without asking
 - effectively, back to circuit-switched
- Higher-level security services \rightarrow more application-layer access via gateways, proxies, ...
- User identity
 - problem is not availability, but rather over-abundance

Scaling

- Scaling is only backbone problem
- Depends on network evolution:
 - continuing addition of AS to flat space \rightarrow deep trouble
 - additional hierarchy

Quality of Service (QoS)

- QoS is meaningless to users
- · care about service availability \rightarrow reliability
- as more and more value depends on network services, can't afford random downtimes

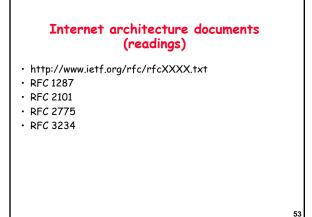
Textbook Internet vs. real Internet

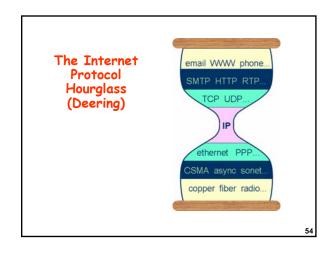
end-to-end (application only in 2 places)	middle boxes (proxies, ALGs,)
permanent interface identifier (IP address)	time-varying (DHCP)
globally unique and routable	network address translation (NAT)
multitude of L2 protocols (ATM, ARCnet, Ethernet, FDDI, modems,)	dominance of Ethernet, but also L2's not designed for networks (1394 Firewire, Fibre Channel, MPEG2,)

Textbook Internet vs. real Internet

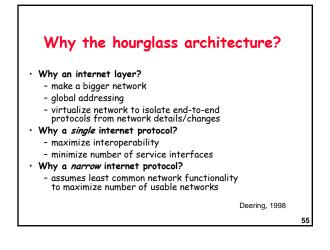
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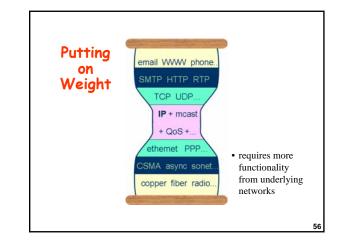
mostly trusted end users	hackers, spammers, con artists, pornographers,
small number of manufacturers, making expensive boxes	Linksys, Dlink, Netgear,, available at Radio Shack
technical users, excited about new technology	grandma, frustrated if email doesn't work
4 layers (link, network, transport, application)	layer splits
transparent network	firewalls, L7 filters, "transparent proxies"

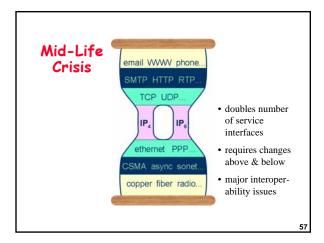


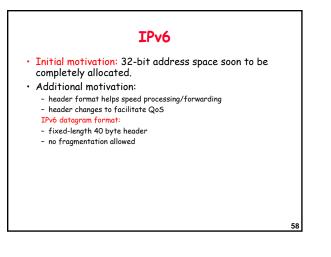


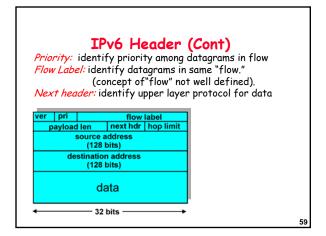
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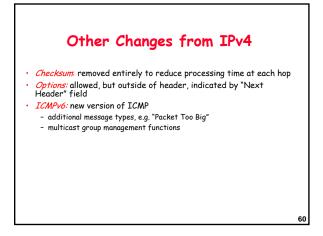


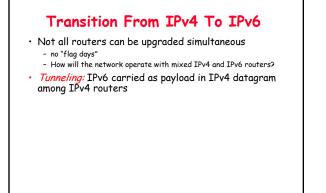




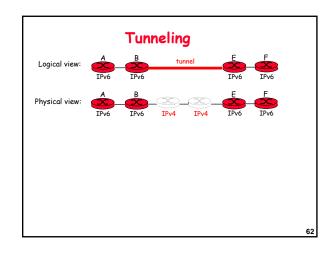


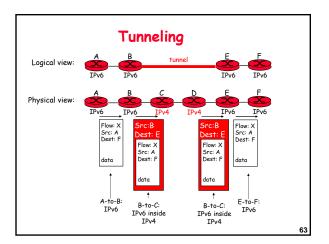


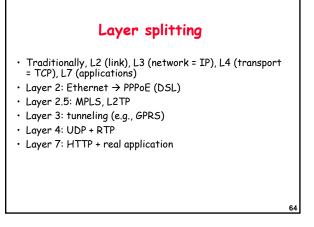




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Layer violations

- + Layers offer abstraction \rightarrow avoid "Internet closed for renovation"
- $\boldsymbol{\cdot}$ Cost of information hiding
- Cost of duplication of information when nothing changes
 fundamental design choice of Internet = difference between circuit and datagram-oriented networks
- Assumption: packets are large and getting larger
- wrong for games and audio
- $\cdot \ \mbox{Cost}$ prohibitive on wireless networks
 - will see: 10 bytes of payloads, 40 bytes of packet header
 - header compression \rightarrow compress into state index on one link

Internet acquires presentation layer All learn about OSI 7-layer model OSI: ASN.1 as common rendering of application data structures used in LDAP and SNMP (and H.323) Internet never really had presentation layer approximations: common encoding (TLV, RFC 822 styles) Now, XML as the design choice by default

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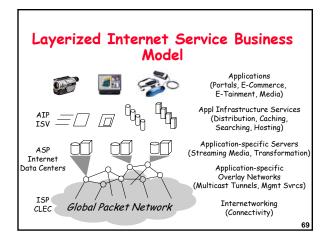
- Originally, meant for data sessions
- Example (not explicit): ftp control connection
- Now, separate data delivery from session setup address and application configuration
 - address and applicati
 deal with mobility
 - will see as RTSP, SIP and H.323

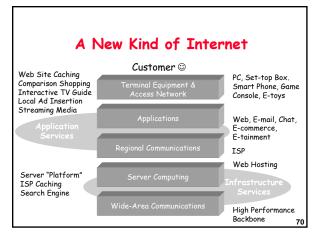
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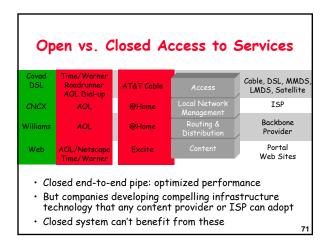
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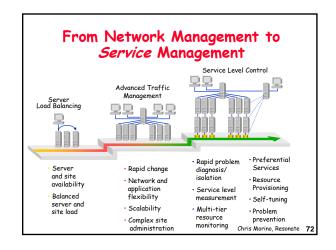
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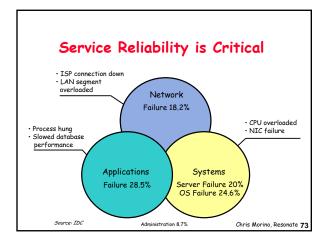
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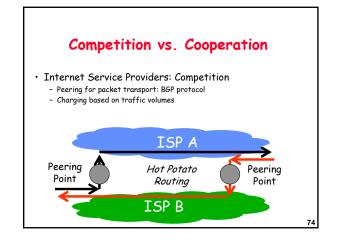






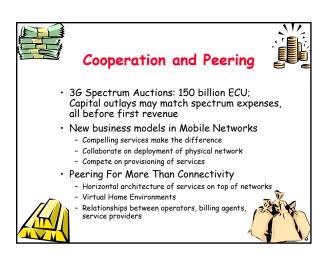






Mobile Internet Might Be Different Than Wired Internet

- Wireless is a smarter pipe - Location-awareness
- UI dictates need for personalization, mediation
- Clear billing authority: it's the access provider - People actually do pay for transport
- Reverse billing allows contact provider to charge for service
 Peering as a necessity
- Operators provide local service
- Roaming agreements provide basis for service peering
- Well understood arrangements for settlements
- New economies driving towards shared network deployment
- Person-to-Person communications is a killer app
- Microsoft's non-monopoly

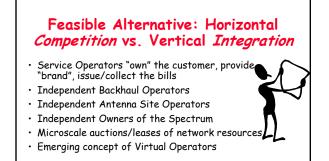


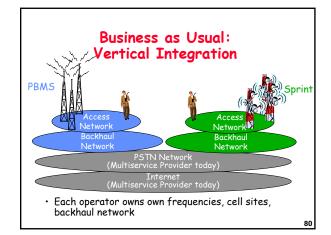
Any Way to Build a Network?

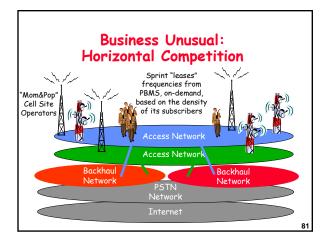
- Partitioning of frequencies independent of actual subscriber density
 - Successful operator oversubscribe resources, while less popular providers retain excess capacity
 - Different flavor of roaming: among collocated/competing service providing
- Duplicate antenna sites
- Serious problem given community resistance
- Redundant backhaul networks - Limited economies of scale

The Case for Horizontal Architectures

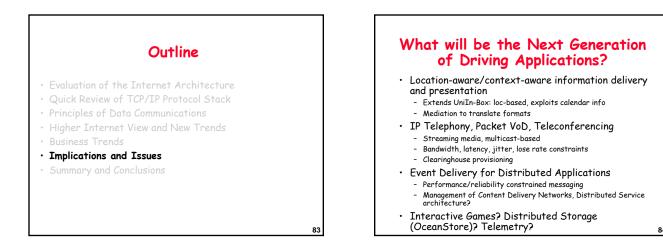
"The new rules for success will be to provide one part of the puzzle and to cooperate with other suppliers to create the complete solutions that customers require. ... [V]ertical integration breaks down when innovation speeds up. The big telecoms firms that will win back investor confidence soonest will be those with the courage to rip apart their monolithic structure along functional layers, to swap size for speed and to embrace rather than fear disruptive technologies." *The Economist Magazine*, 16 December 2000











What Will Be the Next Generation **Operational Environment?**

- Virtual Operators/Service Provider (VOSP) Provide service to end users with no server/network infrastructure of own
 - Independent "Path" providers (e.g., ISPs) and Server providers (e.g., Internet Data Centers)

 - Many-to-many relationship between VOSP and Path/Server Providers
- Confederated Service Provider Service-level peering: sharing of paths and servers to deploy end-to-end service with performance and reliability constraints
- Note: Akamai runs "the world's largest service network" without owning a network!

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Alternative Operational Environments

- Confederation Model
 - Providers share (limited) information about topology, server location, path performance Cooperatively collect internal information and share
- Overlay Model
 - Reverse-engineer topology and intra-cloud performance - Collection done by brokers outside of the cloud
- SLAs, Verification, Maintenance of Trust Relationships different in the two models
- Is there an operational/performance advantage to the Confederation Model?

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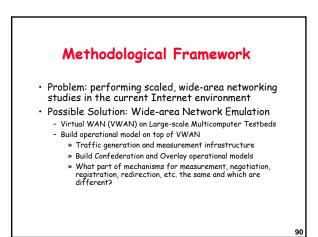
Open Issues/Questions

- Traditional Overlay Networks
 - Server ("Application Level Router") Placement
 - » For scaling, reliability, load balancing, latency
 - » Where? Network topology discovery: WAN Core, Metro/Regional, Access Networks
 - Choice of Inter-Server "Paths"
 - » For server-to-server latency/bandwidth/loss rate
 - » Predictable/verifiable network performance (intra-ISP
 - SLA)
 - Redirection Mechanisms
 - » Random, round-robin, load-informed redirection
 - » Net vs. server as bottleneck

Open Issues/Questions

- · Performance-constrained Service Placement
 - Separation of Service, Server, Service Path » Assume "Server Centers" known, can be "discovered" (how does OceanStore deal with this?), or register with a Service Placement Service (SPS)
 - » How is Service named, described, performance constraints expressed, and registered?
 - » How is app/service-specific performance measured and made known to Service Placement Service?
 - Brokering between Server Centers and Service Creator, Path Provider and Service Creator
- If core network bandwidth becomes infinite and "free", does it matter where services are placed? - Latency reduction vs. economies of centralized management

Emerging Reference Architecture Distributed Application Marshal Resources Constraint Adapt Specification Based on Economic Constraints Service Service Registration Redirection Pricing Path Broker Server Broker Perf Measurement Service Server Registration Advertisement Verify SLAs Registration Path Provider (ISP Cloud) Server Center Provider

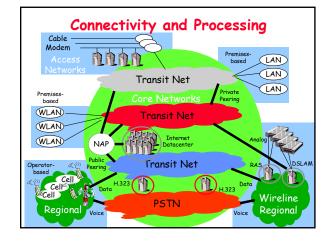


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Challenges for Converged Networks

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- Services spanning access networks, to achieve high performance and manage diversity of end devices
- Not about specific Information Appliances
- Builds on the New Internet: multiple application-specific "overlay" networks, with new kinds of service-level peering
- Pervasive support for services within "intelligent" networks
 - Automatic replication
 - Document routing to caches
 - Compression & mirroring - Data transformation

Managing Edge Versus Core Services

- Wide-area bandwidth efficiency
- Increasing b/w over access networks, but impedance mismatch between core and access nets
- Fast response time (and more predictable)
- Opportunity to untegrate localized content
- Associated with client (actually ISP), not server
- Examples:
- Caching: exploits response time, b/w efficiency, high local b/w
- Filtering: form of local content transformation
- Internet TV: b/w efficiency, high local b/w, predictable response
- Transformation: adapt content for end user/diverse access devices
- Software Rental: sxploits high local b/w - Games, chat rooms,

Yielding a New Research Agenda

- New Definition of "Quality of Service" - Perceived quality depends on services in the network
 - Manage caches, redirection, NOT bandwidth
 - Enable incorporation of localized content
- Bandwidth Issues
 - Tier 1 ISP backbones rapidly moving towards OC 192 (9.6 gbs!) Better interconnection: hops across ASs decreasing over time

 - Emerging broadband access networks: cable, DSL, ...
- End-to-end latency/server load dominate performance · Supporting Old Services in the New Internet
- IP Multicast, DNS,
- Rethinking the End-to-End Principle
- Service/content-level peering, just like routing-level peering
- Secure end-to-end connection compatible with service model?

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